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## Loop Unrolling

## Source:

for $i:=1$ to 100 by 1
A[i] := A[i] + B[i];
endfor

## Transformed Code:

```
    for i \(:=1\) to 100 by 4
        \(\mathrm{A}[\mathrm{i}] \quad:=\mathrm{A}[\mathrm{i}]+\mathrm{B}[\mathrm{i}]\);
        \(\mathrm{A}[i+1]:=\mathrm{A}[i+1]+\mathrm{B}[i+1]\);
        \(\mathrm{A}[\mathrm{i}+2]:=\mathrm{A}[\mathrm{i}+2]+\mathrm{B}[\mathrm{i}+2]\);
        \(\mathrm{A}[i+3]:=\mathrm{A}[i+3]+\mathrm{B}[i+3]\);
    endfor
                                    Benefits:
                            - The overhead of testing and
                        branching is reduced.
                            - This optimization may
                                "enable" other optimizations.
```

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Loop-Invariant Computations
An assignment

$$
\mathbf{x}:=\mathbf{y} \oplus \mathbf{z}
$$

is "Loop-Invariant" if..

- It is in a loop, and
- All definitions of $y$ and $z$ that reach the statement are outside the loop.

We may be able to move the computation into the "preheader".

Step 1: Detect the Loop-Invariant Computations.
Step 2: See if it is okay to move the statement into the pre-header.
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## Detecting Loop-Invariant Computations

## Input:

Loop L (= a set of basic blocks)
U-D Chain information

## Output:

The set of loop-invariant statements.

## Idea:

- Mark some of the statements as "loop-invariant".
- This may allow us to mark even more statements as loop-invariant.
- Remember the order in which theses statements are marked.


## Detecting Loop-Invariant Computations

```
repeat until no new statements are marked...
    Look at each statement in the loop.
    If all its operands are unchanging then
        mark the statement as "loop-invariant".
    An operand is "unchanging" if...
            - It is a constant
            - It has all reaching definitions
                outside of the loop
            - It has exactly one reaching definition
                and that definition has already
                been marked "loop-invariant".
end
```

Remember the order in which statements are marked "loop-invariant."
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## Moving Loop-Invariant Computations

## Consider moving statement

$$
\text { S: } x:=y \oplus z
$$

into the loop's preheader.
The statement must satisfy three conditions.
If it satisfies all conditions, then it can be moved.

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## Condition 3

All uses of " $x$ " in the loop must be reached by ONLY the loop-invariant assignment.


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If all three conditions are satisfied, move the statements into the preheader in the order they were marked Loop-Invariant.



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