CS 311: Computational Structures

Practice Final Exam

The exam is a closed book exam. Answer questions in the space provided.

Student Name (neatly in block letters): _____

All questions from the midterm and quizzes are good study problems. You should familiarize yourself with these questions and their answers. A version of one of the questions in this practice exam will appear in the actual final exam. The actual final exam will also have a pumping lemma question.

1. CFG to PDA

Consider the Context Free Grammar for a language we will call L:

$$egin{array}{ccc} S &
ightarrow & SaSb \ & \mid & \Lambda \end{array}$$

- (a) Give a leftmost derivation for the string *abaabb*.
- (b) Construct a PDA M from the grammar that accepts L. (You may assume a generalized version of the PDA that makes multiple stack moves on a single transition.)
- (c) Show that M accepts the string *abaabb*.
- 2. **Proof by induction**. (20 points) In a earlier homework you showed by induction, that for one representation of the natural numbers, and for one inductive definition of addition, that addition was associative, that is i + (j + k) = (i + j) + k. In this problem we use the same representation of the natural numbers and addition, but instead show that the following property holds: for all x the following holds x + (Sy) = S(x + y). Recall the definitions for the natural numbers and addition (infix +).
 - Z is a natural number (zero)
 - if n is a natural number, then S n is a natural number (the successor function).
 - Nothing else is a natural number
 - (1) Z + n = n(2) (S m) + n = S(m + n)

Prove by structural induction that for all x the following property holds: x + (Sy) = S(x + y).

Be sure and structure your proof properly, breaking it into appropriate cases, and clearly stating the induction hypothesis where appropriate.

3. Turing computable.

- State the Church-Turing Thesis.
- List three computational systems that are Turing complete.
- Describe a string acceptance problem that is computable by a Turing machine, but not by a context free grammar.
- 4. **Traces.** Correctness (acceptance) of many computational systems can be defined in terms of the existence of a trace.
 - What important property do traces have?
 - List three of the systems described in the Outline notes, and describe what their traces look like.

- Consider the context free grammars. Outline what a trace for a context free grammar might look like.
- Define a simple CFG, and provide a trace in the form you described above for a short string.

5. Computability.

- What is a universal machine?
- What important property does a universal machine have?
- Give an example of a universal machine that we saw in the lecture notes.
- Are all total functions primitive recursive? If not give an example.
- What is the halting problem?