

# CS 311: Computational Structures

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## 4 GNFA Example in Detail

This note illustrates the application of the construction in Sipser’s proof of Lemma 1.60 to the modulo 3 counter DFA.

In lecture we have shown the DFA that recognizes binary numbers modulo 3.  $M = (\{0, 1, 2\}, \{0, 1\}, \delta, 0, \{2\})$ , where  $\delta$  is given by the table:

$$\begin{aligned} \delta(0,0) &= 0 \\ \delta(0,1) &= 1 \\ \delta(1,0) &= 2 \\ \delta(1,1) &= 0 \\ \delta(2,0) &= 1 \\ \delta(2,1) &= 2 \end{aligned}$$

Convert this to a GNFA by adding states  $s$  and  $a$ , and labeling all transitions with regular expressions. The GNFA can be seen by this table, where each row is the “from” state and each column is the “to” state:

$\delta$	0	1	2	$a$
$s$	$\epsilon$	$\emptyset$	$\emptyset$	$\emptyset$
0	0	1	$\emptyset$	$\emptyset$
1	1	$\emptyset$	0	$\emptyset$
2	$\emptyset$	0	1	$\epsilon$

To “rip” state 0 out of this machine, we calculate the following  $\delta'$ :

$\delta$	1	2	$a$
$s$	$\delta_{s0}\delta_{00}^*\delta_{01} \cup \delta_{s1}$	$\delta_{s0}\delta_{00}^*\delta_{02} \cup \delta_{s2}$	$\delta_{s0}\delta_{00}^*\delta_{0a} \cup \delta_{sa}$
1	$\delta_{10}\delta_{00}^*\delta_{01} \cup \delta_{11}$	$\delta_{10}\delta_{00}^*\delta_{02} \cup \delta_{12}$	$\delta_{10}\delta_{00}^*\delta_{0a} \cup \delta_{1a}$
2	$\delta_{20}\delta_{00}^*\delta_{01} \cup \delta_{21}$	$\delta_{20}\delta_{00}^*\delta_{02} \cup \delta_{22}$	$\delta_{20}\delta_{00}^*\delta_{0a} \cup \delta_{2a}$

$\delta$	1	2	$a$
$s$	$\epsilon 0^* 1 \cup \emptyset$	$\epsilon 0^* \emptyset \cup \emptyset$	$\epsilon 0^* \emptyset \cup \emptyset$
1	$10^* 1 \cup \emptyset$	$10^* \emptyset \cup 0$	$10^* \emptyset \cup \emptyset$
2	$\emptyset 0^* 1 \cup 0$	$\emptyset 0^* \emptyset \cup 1$	$\emptyset 0^* \emptyset \cup \epsilon$

$\delta$	1	2	$a$
$s$	$0^*1$	$\emptyset$	$\emptyset$
1	$10^*1$	0	$\emptyset$
2	0	1	$\epsilon$

Next, rip state 2 out:

$\delta$	1	$a$
$s$	$\delta_{s2}\delta_{22}^*\delta_{21} \cup \delta_{s1}$	$\delta_{s2}\delta_{22}^*\delta_{2a} \cup \delta_{sa}$
1	$\delta_{12}\delta_{22}^*\delta_{21} \cup \delta_{11}$	$\delta_{12}\delta_{22}^*\delta_{2a} \cup \delta_{1a}$

$\delta$	1	$a$
$s$	$\emptyset 1^* 0 \cup 0^* 1$	$\emptyset 1^* \epsilon \cup \emptyset$
1	$0 1^* 0 \cup 1 0^* 1$	$0 1^* \epsilon \cup \emptyset$

$\delta$	1	$a$
$s$	$0^*1$	$\emptyset$
1	$0 1^* 0 \cup 1 0^* 1$	$0 1^*$

Finally, we rip out state 1, leaving the single transition:

$$\delta_{s1}\delta_{11}^*\delta_{1a} \cup \delta_{sa}$$

Which becomes:

$$0^*1(0 1^* 0 \cup 1 0^* 1)^* 0 1^* \cup \emptyset$$

Or simply:

$$0^*1(0 1^* 0 \cup 1 0^* 1)^* 0 1^*$$

The result is a regular expression generating the set of binary numbers that are congruent to 2 modulo 3.

**Exercise 4.1** What if we ripped the states in a different order?